



# Algebraic Effects on the JVM

**Jonathan Immanuel Brachthäuser**  
University of Tübingen, Germany

**Dagstuhl Seminar 18172**  
Algebraic Effect Handlers go Mainstream



[github.com/b-studios/scala-effekt](https://github.com/b-studios/scala-effekt)

# We developed algebraic effect libraries for



**E** Scala Effekt



**E** JVM Effekt

# Part I

## Effect

### Handlers as a Library for Scala

# Key Specs: Scala Effekt

- shallow embedding vs. deep embedding of handlers
- "capability passing style"
- shallow handlers vs. deep handlers
- user defined effects ✓
- dynamic effect instances ✓
- modular and extensible effect signatures and handlers ✓
- safety (capabilities can leak) ✗
- user programs are written in direct style ✗
- performance: still (orders of magnitude) slower than primitive effects ✗

# How to **use** Scala Effekt

# Example: Drunk Coin Flipping

```
val drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads   ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

```
val drunkFlip: String using Amb and Exc = for {  
    caught ← flip()  
    heads ← if (caught) flip() else raise("too drunk")  
} yield if (heads) "heads" else "tails"
```

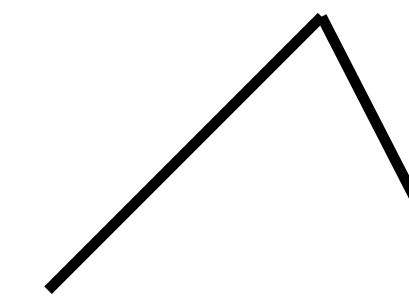
## Effect Operations

Semantics of the operations is left open

## Effect Signatures

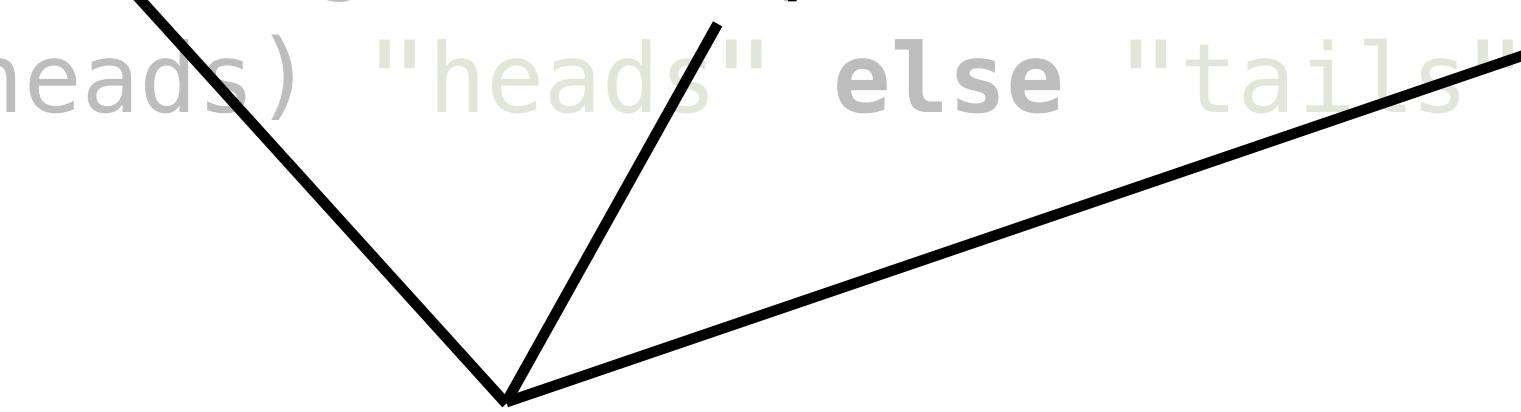
Group effect operations in one type

```
val drunkFlip: String using Amb and Exc = for {  
    caught ← flip()  
    heads ← if (caught) flip() else raise("too drunk")  
} yield if (heads) "heads" else "tails"
```



## Effect Operations

Semantics of the operations is left open



# ≡ Effekt

```
val drunkFlip: String using Amb and Exc = for {  
    caught ← flip()  
    heads ← if (caught) flip() else raise("too drunk")  
} yield if (heads) "heads" else "tails"
```

```
AmbList { ExcOption { drunkFlip } }
```

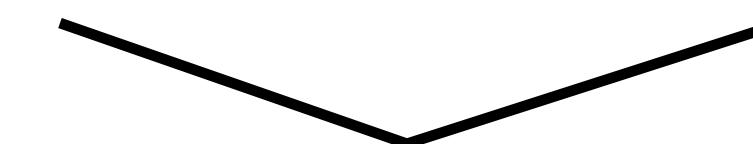
# ≡ Effekt

```
val drunkFlip: String using Amb and Exc = for {  
    caught ← flip()  
    heads ← if (caught) flip() else raise("too drunk")  
} yield if (heads) "heads" else "tails"
```

08

b-studios.de/scala-effekt

AmbList { ExcOption { drunkFlip } }



## Effect Handlers

Provide semantics to effect operations

# ≡ Effekt

```
val drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

```
AmbList { ExcOption { drunkFlip } }
```

# ≡ Effekt

```
val drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads  ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

```
AmbList { ExcOption { drunkFlip } }
> List(Some("heads"), Some("tails"), None)
```

# Effekt

```
val drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads  ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

```
AmbList { ExcOption { drunkFlip } }
> List(Some("heads"), Some("tails"), None)
```

```
ExcOption { AmbList { drunkFlip } }
```

# Effekt

```
val drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads   ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

```
AmbList { ExcOption { drunkFlip } }
> List(Some("heads"), Some("tails"), None)
```

```
ExcOption { AmbList { drunkFlip } }
> None
```

# The role of implicit

# Design Decisions - Scala Effekt

We were faced with the following three design questions:

# Design Decisions - Scala Effekt

We were faced with the following three design questions:

1. how to **capture the context** of effect operations?  
    ⇒ monadic implementation for multi-prompt delimited continuations

# Design Decisions - Sca

We were faced with the following three

1. how to **capture the context** of effects  
⇒ monadic implementation for monads

To appear in *J. Functional Programming*

1

## *A Monadic Framework for Delimited Continuations*

R. Kent Dywig  
Indiana University      Simon Peyton Jones  
Microsoft Research      Amr Sabry\*

---

### Abstract

Delimited continuations are more expressive than traditional abortive continuations and they apparently require a framework beyond traditional continuation-passing style (CPS). We show that this is not the case: standard CPS is sufficient to explain the common control operators for delimited continuations. We demonstrate this fact and present an implementation as a Scheme library. We then investigate a typed account of delimited continuations that makes explicit where control effects can occur. This results in a monadic framework for typed and encapsulated delimited continuations, which we design and implement as a Haskell library.

---

### 1 Introduction

Continuation-passing style (CPS) and its generalisation to monadic style are the standard mathematical frameworks for understanding (and sometimes implementing) control operators. Starting in the late eighties a new class of control operators were introduced that apparently went “beyond continuations” (Pelleisen *et al.*, 1987b; Felleisen, 1988; Hieb & Dywig, 1990; Danvy & Filinski, 1990) and “beyond monads” (Wadler, 1994). These control operators permit the manipulation of *delimited continuations* that represent only part of the remainder of a computation, and they also support the *composition* of continuations, even though such operations are not directly supported by standard continuation models (Strachey & Wadsworth, 1974). Delimited continuations are also referred to as *subcontinuations* (Hieb *et al.*, 1994), since they represent the remainder of a subcomputation rather than of a computation as a whole. We use the terms interchangeably in this article.

Without the unifying frameworks of CPS or monads, it is difficult to understand, compare, implement, and reason about the various control operators for delimited continuations, their typing properties, and logical foundations. In this article, we design such a unifying framework based on continuation semantics, then generalise it to a typed monadic semantics. We illustrate this framework with a basic set of

\* Supported by National Science Foundation grant number CCR-0196063, by a Visiting Researcher position at Microsoft Research, Cambridge, U.K., and by a Visiting Professor position at the University of Genova, Italy.

# Design Decisions - Scala Effekt

1. how to **capture the context** of effect operations?  
    ⇒ monadic implementation for multi-prompt delimited continuations
  
2. how should **effect handlers provide semantics** for effect operations?  
    ⇒ shallow embedding of effect handlers

# Design Decisions - Scala Effekt

We were faced with the following three design questions:

1. how to **capture the context** of effect operations?  
    ⇒ monadic implementation for multi-prompt delimited continuations
2. how should **effect handlers provide semantics** for effect operations?  
    ⇒ shallow embedding of effect handlers
3. how to establish an **effect typing discipline**?  
    ⇒ capability passing style

# Design Decisions - Scala Effekt

We were faced with the following three design questions:

1. how to **capture the context** of effect operations?  
⇒ monadic implementation for multi-prompt delimited continuations
2. how should **effect handlers provide semantics** for effect operations?  
⇒ shallow embedding of effect handlers
3. how to establish an **effect typing discipline**?  
⇒ capability passing style

For all three answers **implicit function types** turned out to be a perfect fit!

# Implicit Function Types

# Implicit Function Types

```
def drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads   ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

# Implicit Function Types

```
def drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads   ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

Library defined type aliases

```
type using [A, E] = implicit Cap[E] ⇒ Control[A]
```

# Implicit Function Types

```
def drunkFlip: String using Amb and Exc = for {
    caught ← flip()
    heads   ← if (caught) flip() else raise("too drunk")
} yield if (heads) "heads" else "tails"
```

Library defined type aliases

```
type using [A, E] = implicit Cap[E] ⇒ Control[A]
```

```
type and [A, E]   = implicit Cap[E] ⇒ A
```

# Making Capability Passing Explicit

Explicitly desugaring implicit function types gives:

013

```
def drunkFlip(amb: Cap[Amb], exc: Cap[Exc]): Control[String] = for {
    caught ← amb.handler.flip()
    heads   ← if (caught) amb.handler.flip()
              else exc.handler.raise("too drunk")
} yield if (heads) "heads" else "tails"
```

# Making Capability Passing Explicit

Explicitly desugaring implicit function types gives:

```
def drunkFlip(amb: Cap[Amb], exc: Cap[Exc]): Control[String] = for {
    caught ← amb.handler.flip()
    heads ← if (caught) amb.handler.flip()
             else exc.handler.raise("too drunk")
} yield if (heads) "heads" else "tails"
```

# Making Capability Passing Explicit

Explicitly desugaring implicit function types gives:

```
def drunkFlip(amb: Cap[Amb], exc: Cap[Exc]): Control[String] = for {
    caught ← amb.handler.flip()
    heads ← if (caught) amb.handler.flip()
             else exc.handler.raise("too drunk")
} yield if (heads) "heads" else "tails"
```

Handler create capabilities:

```
AmbList { amb ⇒ ExcOption { exc ⇒ drunkFlip(amb, exc) } }
```

# Capabilities in Effekt

Capabilities  $\text{Cap}[E]$  encapsulate three different things:

# Capabilities in Effekt

Capabilities  $\text{Cap}[E]$  encapsulate three different things:

1. they contain a unique prompt marker that **delimits the scope** of the continuation to be captured.

# Capabilities in Effekt

Capabilities  $\text{Cap}[E]$  encapsulate three different things:

1. they contain a unique prompt marker that **delimits the scope** of the continuation to be captured.
2. they contain the **effect handler implementation** to be passed down (shallow embedding of handlers).

# Capabilities in Effekt

Capabilities  $\text{Cap}[E]$  encapsulate three different things:

1. they contain a unique prompt marker that **delimits the scope** of the continuation to be captured.
2. they contain the **effect handler implementation** to be passed down (shallow embedding of handlers).
3. they entitle the holder of the capability to **invoke effectful operations** specified in effect signature  $E$  (effect typing discipline).



# shallow embedding

of effect handlers

# Calling an Effect Operation

# Calling an Effect Operation

We can think of effect operations as uninterpreted constructors of an effect-language. An effectful program then could be represented as a tree of operations:

$$\begin{aligned} 0p_1(\textit{args}\dots, \textit{res}_1 \Rightarrow \\ 0p_2(\textit{args}\dots, \textit{res}_2 \Rightarrow \\ \dots \\ \textit{Pure}(\textit{value}))) \end{aligned}$$

# Calling an Effect Operation

We can think of effect operations as uninterpreted constructors of an effect-language. An effectful program then could be represented as a tree of operations:

$$\begin{aligned} 0p_1(args\ldots, res_1 \Rightarrow \\ 0p_2(args\ldots, res_2 \Rightarrow \\ \cdots \\ Pure(value))) \end{aligned}$$

we can write a recursive, pattern matching recursive interpreter to provide semantics to effectful operations.

In PL terms: a *deep embedding* of effect operations.

# Shallow Embedding of Effect Handlers

In Scala Effekt, effect operations are immediately called on effect handlers.

*Schematically:*

```
handler.op1(args..., res1 =>  
  handler.op2(args..., res2 =>  
    ...))
```

# Shallow Embedding of Effect Handlers

In Scala Effekt, effect operations are immediately called on effect handlers.

*Schematically:*

```
handler.op1(args..., res1 =>  
  handler.op2(args..., res2 =>  
    ...))
```

## Technical Insights

# Shallow Embedding of Effect Handlers

In Scala Effekt, effect operations are immediately called on effect handlers.

*Schematically:*

```
handler.op1(args..., res1 =>  
  handler.op2(args..., res2 =>  
    ...))
```

## Technical Insights

- (a) Shallow embedding of effect handlers simplifies typing – no GADTs are necessary!

# Shallow Embedding of Effect Handlers

In Scala Effekt, effect operations are immediately called on effect handlers.

*Schematically:*

```
handler.op1(args..., res1 =>  
  handler.op2(args..., res2 =>  
    ...))
```

## Technical Insights

- (a) Shallow embedding of effect handlers simplifies typing – no GADTs are necessary!
- (b) Pattern matching is replaced by dynamic dispatch – benefits performance on the JVM.

# Shallow Embedding of Effect Handlers

In Scala Effekt, effect operations are immediately called on effect handlers.

*Schematically:*

```
handler.op1(args..., res1 =>  
  handler.op2(args..., res2 =>  
    ...))
```

## Technical Insights

- (a) Shallow embedding of effect handlers simplifies typing – no GADTs are necessary!
- (b) Pattern matching is replaced by dynamic dispatch – benefits performance on the JVM.
- (c) Direct call to corresponding handler – no need to lookup handler.

# Part II

## Algebraic Effects as Libraries for Java / JVM

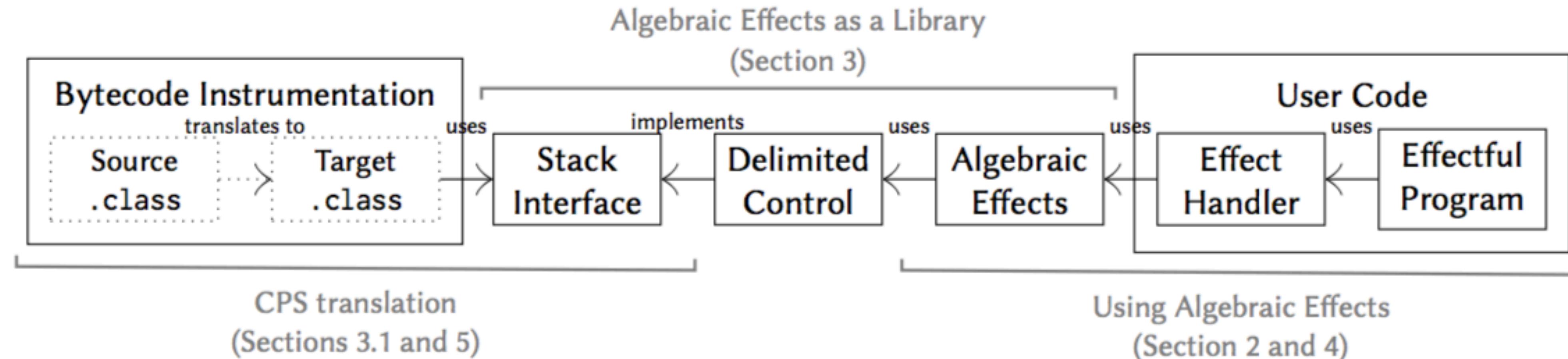
# Key Specs: JVM / Java Effekt

- shallow embedding vs. deep embedding of handlers
- "handler passing style"
- shallow handlers vs. deep handlers
- user defined effects ✓
- dynamic effect instances ✓
- modular and extensible effect signatures and handlers (✓)
- safety (capabilities can leak) ✗
- user programs are written in direct style ✓
- performance: competitive with JVM continuation libraries ✓

# Key Specs: JVM / Java Effekt

- shallow embedding vs. deep embedding of handlers
- "handler passing style"
- shallow handlers vs. deep handlers
- user defined effects ✓
- dynamic effect instances ✓
- modular and extensible effect signatures and handlers (✓)
- safety (capabilities can leak) ✗
- user programs are written in direct style ✓
- performance: competitive with JVM continuation libraries ✓

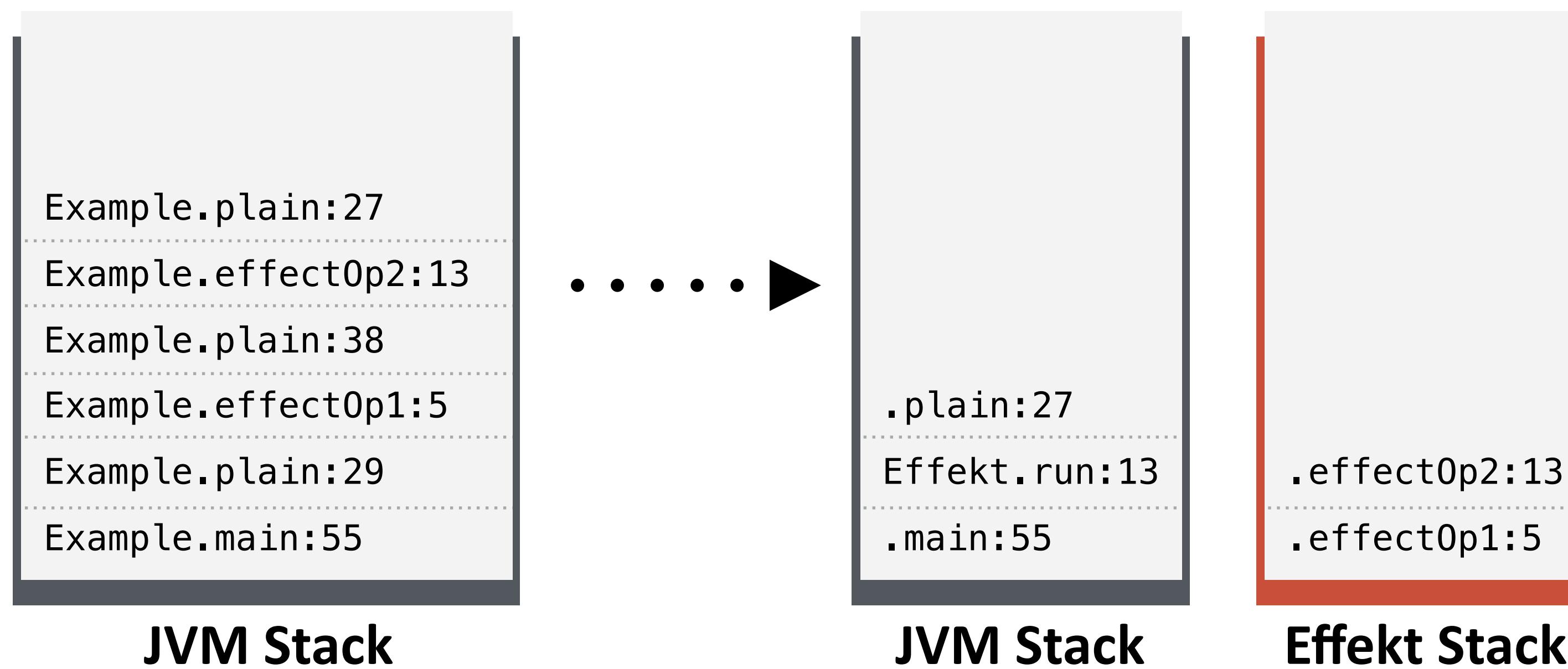
# Overview of JVM Effekt



- Programs are written in direct style, but CPS translated via **bytecode transformation**
- Translated programs use a separate Stack interface for effectful frames
- Delimited control is implemented as a library, implementing the Stack interface
- We redesigned the algebraic effects library to **only require simple generics**
- **Restriction:** We only transform the terms, not types / signatures

# Replacing the JVM Stack

For effectful methods, we maintain our own custom stack, which allows us to manipulate it (searching, slicing, copying).



# Example: Drunk Coin Flipping

```
String drunkFlip(Amb amb, Exc exc) throws Effects {
    if (amb.flip()) {
        return exc.raise("too drunk");
    } else {
        return amb.flip() ? "heads" : "tails";
    }
}
```

# Example: Drunk Coin Flipping

```
interface Amb { boolean flip() throws Effects; }
interface Exc { <A> A raise(String msg) throws Effects; }
```

```
String drunkFlip(Amb amb, Exc exc) throws Effects {
    if (amb.flip()) {
        return exc.raise("too drunk");
    } else {
        return amb.flip() ? "heads" : "tails";
    }
}
```

# Handling Effects

```
class AmbList<R> extends Handler<R, List<R>> implements Amb {
    List<R> pure(R r) { return Lists.singleton(r); }
    boolean flip() throws Effects {
        return use(k -> Lists.concat(k.resume(true), k.resume(false)));
    }
}

handle(new AmbList<Optional<String>>(), amb ->
    handle(new Maybe<String>(), exc -> drunkFlip(amb, exc)))

> [Optional["heads"], Optional["tails"], Optional.empty]
```

# Stateful / Parametrized Handlers

```
interface Reader<In> { In read() throws Effects; }

class StringReader<R> extends Handler<R, R> implements Reader<Char> {

    final String input;
    int pos = 0;

    Char read() throws Effects { return input.charAt(pos++) }
}
```

# Stateful / Parametrized Handlers

```
interface Reader<In> { In read() throws Effects; }

class StringReader<R> extends Handler<R, R> implements Reader<Char>,
    Stateful<Integer> {
    final String input;
    int pos = 0;

    Char read() throws Effects { return input.charAt(pos++) }

    Integer exportState() { return pos; }
    void importState(Integer n) { pos = n; }
}
```

# Design Decisions

- **Effectful methods** are marked with a special, checked exception Effects
- **Effect signatures** are interfaces that contain effectful methods
- **Effect handlers** are implementations of those interfaces.
- Users need to manually follow the **capability passing style**.
- Effect handlers can extend the library class Handler to **capture the continuation** (but don't need to).
- We use the handler instances as **prompt markers**.

# Bytecode Transformation Example (CPS)

```
String drunkFlip(Amb amb, Exc exc) throws Effects {  
    Effekt.push(() -> drunkFlip1(amb, exc));  
    amb.flip();  
    return null;  
}
```

# Bytecode Transformation Example (CPS)

```
String drunkFlip(Amb amb, Exc exc) throws Effects {
    Effekt.push(() -> drunkFlip1(amb, exc));
    amb.flip();
    return null;
}

void drunkFlip1(Amb amb, Exc exc) throws Effects {
    boolean caught = Effekt.result();
    if (Effekt.result()) { exc.raise("too drunk"); }
    else {
        Effekt.push(() -> drunkFlip2(amb, exc, caught));
        amb.flip();
    }
}
```

# Bytecode Transformation Example (CPS)

```
String drunkFlip(Amb amb, Exc exc) throws Effects {
    Effekt.push(() -> drunkFlip1(amb, exc));
    amb.flip();
    return null;
}

void drunkFlip1(Amb amb, Exc exc) throws Effects {
    boolean caught = Effekt.result();
    if (Effekt.result()) { exc.raise("too drunk"); }
    else {
        Effekt.push(() -> drunkFlip2(amb, exc, caught));
        amb.flip();
    }
}

void drunkFlip2(Amb amb, Exc exc, boolean caught) throws Effects {
    Effekt.returnWith(Effekt.result() ? "heads" : "tails");
}
```

# Alternative Transformations

## CPS

```
Effekt.push(() -> drunkFlip1(amb, exc));  
amb.flip();  
return DUMMY;
```

# Alternative Transformations

## CPS

```
Effekt.push(() -> drunkFlip1(amb, exc));  
amb.flip();  
return DUMMY;
```

## Gen. Stack Inspection / Bubble Sem.

```
Effekt.beforeCall();  
amb.flip();  
if (Effekt.isImpure()) {  
    Effekt.push(() -> drunkFlip1(amb, exc));  
    return DUMMY;  
}
```

# Alternative Transformations

## CPS

```
Effekt.push(() -> drunkFlip1(amb, exc));  
amb.flip();  
return DUMMY;
```

## Gen. Stack Inspection / Bubble Sem.

```
Effekt.beforeCall();  
amb.flip();  
if (Effekt.isImpure()) {  
    Effekt.push(() -> drunkFlip1(amb, exc));  
    return DUMMY;  
}
```

- all effect calls are tail calls
- cont. is constructed eagerly and immediately available
- unnecessary push/pop/enter cycles
- full reification of the stack

# Alternative Transformations

## CPS

```
Effekt.push(() -> drunkFlip1(amb, exc));  
amb.flip();  
return DUMMY;
```

## Gen. Stack Inspection / Bubble Sem.

```
Effekt.beforeCall();  
amb.flip();  
if (Effekt.isImpure()) {  
    Effekt.push(() -> drunkFlip1(amb, exc));  
    return DUMMY;  
}
```

- all effect calls are tail calls
- cont. is constructed eagerly and immediately available
- unnecessary push/pop/enter cycles
- full reification of the stack

- two ways to leave a method, distinguished by a flag
- cont. is constructed on demand
- reduced overhead for pure code
- prompt markers are trampolines

# Alternative Transformations (Performance)

<b>Benchmark</b>	Time in ms (Confidence Interval)					
	<b>Baseline</b>	<b>EFFEKT</b>	<b>EFFEKT<sub>opt</sub></b>	<b>Coroutines</b>	<b>Quasar</b>	<b>JavaFlow</b>
Stateloop 1M	1.61 ±0.09	29.76 ±2.57	1.91 ±0.04	5.52 ±0.35	69.02 ±2.59	14.82 ±0.48
RecursiveOnce 1K	0.01 ±0.0	0.69 ±0.22	0.34 ±0.01	0.07 ±0.0	0.23 ±0.03	8.18 ±0.19
RecursiveMany 1K	0.01 ±0.0	1.05 ±0.38	0.4 ±0.07	10.29 ±1.41	68.07 ±2.07	3363.74 ±23.46
Skynet 1M	2.74 ±0.03	171.34 ±5.55	62.13 ±3.87	35.19 ±2.51	762.1 ±155.95	1277.51 ±54.18
SkynetSuspend 1M	2.74 ±0.03	414.56 ±9.2	147.4 ±5.44	50.46 ±2.95	1113.15 ±112.78	7198.72 ±122.56

# Alternative Transformations (Performance)

<b>Benchmark</b>	Time in ms (Confidence Interval)					
	<b>Baseline</b>	<b>EFFEKT</b>	<b>EFFEKT<sub>opt</sub></b>	<b>Coroutines</b>	<b>Quasar</b>	<b>JavaFlow</b>
Stateloop 1M	1.61 ±0.09	29.76 ±2.57	1.91 ±0.04	5.52 ±0.35	69.02 ±2.59	14.82 ±0.48
RecursiveOnce 1K	0.01 ±0.0	0.69 ±0.22	0.34 ±0.01	0.07 ±0.0	0.23 ±0.03	8.18 ±0.19
RecursiveMany 1K	0.01 ±0.0	1.05 ±0.38	0.4 ±0.07	10.29 ±1.41	68.07 ±2.07	3363.74 ±23.46
Skynet 1M	2.74 ±0.03	171.34 ±5.55	62.13 ±3.87	35.19 ±2.51	762.1 ±155.95	1277.51 ±54.18
SkynetSuspend 1M	2.74 ±0.03	414.56 ±9.2	147.4 ±5.44	50.46 ±2.95	1113.15 ±112.78	7198.72 ±122.56

<b>Benchmark</b>	<b>EFFEKT</b>	<b>EFFEKT<sub>opt</sub></b>	<b>Scala Effekt</b>	<b>Scala Eff</b>
Countdown 10K	3.35 ±0.07	2.47 ±0.12	6.07 ±0.32	34.39 ±2.59
Countdown8 1K	1.31 ±0.39	1.77 ±0.1	2.31 ±0.12	36.92 ±3.0
NQueens (10)	19.5 ±0.38	16.09 ±0.19	40.95 ±0.54	49.89 ±2.17

# Alternative Transformations (Performance)

<b>Benchmark</b>	Time in ms (Confidence Interval)					
	<b>Baseline</b>	<b>EFFEKT</b>	<b>EFFEKT<sub>opt</sub></b>	<b>Coroutines</b>	<b>Quasar</b>	<b>JavaFlow</b>
Stateloop 1M	1.61 ±0.09	29.76 ±2.57	1.91 ±0.04	5.52 ±0.35	69.02 ±2.59	14.82 ±0.48
RecursiveOnce 1K	0.01 ±0.0	0.69 ±0.22	0.34 ±0.01	0.07 ±0.0	0.23 ±0.03	8.18 ±0.19
RecursiveMany 1K	0.01 ±0.0	1.05 ±0.38	0.4 ±0.07	10.29 ±1.41	68.07 ±2.07	3363.74 ±23.46
Skynet 1M	2.74 ±0.03	171.34 ±5.55	62.13 ±3.87	35.19 ±2.51	762.1 ±155.95	1277.51 ±54.18
SkynetSuspend 1M	2.74 ±0.03	414.56 ±9.2	147.4 ±5.44	50.46 ±2.95	1113.15 ±112.78	7198.72 ±122.56

<b>Benchmark</b>	<b>EFFEKT</b>	<b>EFFEKT<sub>opt</sub></b>	<b>Scala Effekt</b>	<b>Scala Eff</b>
Countdown 10K	3.35 ±0.07	2.47 ±0.12	6.07 ±0.32	34.39 ±2.59
Countdown8 1K	1.31 ±0.39	1.77 ±0.1	2.31 ±0.12	36.92 ±3.0
NQueens (10)	19.5 ±0.38	16.09 ±0.19	40.95 ±0.54	49.89 ±2.17

- Coroutines (<https://github.com/offbynull/coroutines>)
- Quasar (<http://docs.paralleluniverse.co/quasar>)
- Javaflow (<https://github.com/vsilaev/tascalate-javaflow>)
- Eff (<https://github.com/atnos-org/eff>)

# Part III

## Even More Extensible Effects

# The (Effect) Expression Problem

*Original Expression Problem*

Variant of a Datatype

*vs.*

*Effect Expression Problem*

Effect Operation

*vs.*

(Recursive) Operation

Handler Implementation

# The (Effect) Expression Problem

<i>Original Expression Problem</i>	<i>vs.</i>	<i>Effect Expression Problem</i>
Variant of a Datatype	<i>vs.</i>	Effect Operation
(Recursive) Operation	<i>vs.</i>	Handler Implementation

We rephrase the expression problem in context of algebraic effects as:

*Modularly being able to*

- implement new handlers for an effect signature.*
- add new effect operations to an existing effect signature*

# Extensibility supported by Effekt

# Extensibility supported by Effekt

a) *implement new handlers for an effect signature.*

```
trait ExcOption[R] extends Exc with Handler[R, Option[R]] { ... }
```

```
trait ExcEither[R] extends Exc with Handler[R, Either[String, R]] { ... }
```

# Extensibility supported by Effekt

a) *implement new handlers for an effect signature.*

```
trait ExcOption[R] extends Exc with Handler[R, Option[R]] { ... }
```

```
trait ExcEither[R] extends Exc with Handler[R, Either[String, R]] { ... }
```

b) *add new effect operations ...*

- ... by adding a new signature (like Amb and Exc)

- ... by adding operations to a signature

# Extensibility supported by Effekt

a) *implement new handlers for an effect signature.*

```
trait ExcOption[R] extends Exc with Handler[R, Option[R]] { ... }
```

```
trait ExcEither[R] extends Exc with Handler[R, Either[String, R]] { ... }
```

b) *add new effect operations ...*

... by adding a new signature (like Amb and Exc)

... by adding operations to a signature

```
trait AmbChoose extends Amb { def choose[A](choices: List[A]): Op[A] }
```

```
trait AmbChooseList[R] extends AmbChoose with AmbList[R] {
```

```
  def choose[A](choices: List[A]): Op[A] = ...
```

```
}
```

# Extensibility supported by Effekt (2)

Handling two effects with one handler:

```
trait ExcList[R] extends Exc with Handler[R, List[R]] {
    def raise[A](msg: String): Op[A] = resume => pure(List.empty)
}

trait ExcAmbList[R] extends ExcList[R] with AmbList[R] {}

ExcAmbList { drunkFlip }

> List("heads", "tails")
```

# Extensibility supported by Effekt (2)

Handling two effects with one handler:

```
trait ExcList[R] extends Exc with Handler[R, List[R]] {  
    def raise[A](msg: String): Op[A] = resume ⇒ pure(List.empty)  
}  
  
trait ExcAmbList[R] extends ExcList[R] with AmbList[R] {}  
  
ExcAmbList { drunkFlip }  
    > List("heads", "tails")
```

Desugares to:

```
ExcAmbList { both ⇒ drunkFlip(both, both) }
```

# Part IV

## Effect Typing and OO: A Problem Statement

# Object Oriented Programming

The mantra of OOP:

- **subtyping** and Liskov's substitution principle
- **hiding implementation details** behind interfaces
- implementation is **existentially** hidden
- Information hiding happens on the granularity of **single objects**

# Subtyping & Information Hiding

```
trait Person {  
    def greet(): Unit  
}
```

```
trait IOPerson extends Person {  
    def greet(): Unit using Console  
}
```

```
trait AlertPerson extends Person {  
    def greet(): Unit using GUI  
}
```

# Subtyping & Information Hiding

```
trait Person {  
    def greet(): Unit  
}
```

```
trait IOPerson extends Person {  
    def greet(): Unit using Console  
}
```

```
trait AlertPerson extends Person {  
    def greet(): Unit using GUI  
}
```

Not a subtype!

# Solution Attempt 1

```
trait Person[E] {  
    def greet(): Unit using E  
}
```

Users of Person now also need to be effect polymorphic!

```
def user[E](p: Person[E]): Int using E
```

```
trait IOPerson extends Person[Console]
```

```
trait AlertPerson extends Person[GUI]
```

# Solution Attempt 1

```
trait Person {  
    type E  
    def greet(): Unit using E  
}
```

Effect types are now path dependent:  
**def user(p: Person): Int using p.E**

```
trait IOPerson extends Person {  
    type E = Console  
}
```

*Only works for stable  
values!*

```
trait AlertPerson extends Person {  
    type E = GUI  
}
```

# Solution Attempt 2

```
trait Person {  
    def greet(): Control[Unit]  
}
```

The effect now is truly hidden  
**def user(p: Person): Control[Int]**

```
trait IOPerson extends Person {  
    implicit val console: Cap[Console]  
}
```

```
trait AlertPerson extends Person {  
    implicit val gui: Cap[GUI]  
}
```

# Capability Safety

```
def leaking(implicit amb: Cap[Amb]): Control[String] = {  
    pure("hello world")  
}
```

# Capability Safety

```
var c: Cap[Amb] = null
def leaking(implicit amb: Cap[Amb]): Control[String] = {
    c = amb;
    pure("hello world")
}
```

# 1 Effekt

## Capability Safety

```
var c: Cap[Amb] = null
def leaking(implicit amb: Cap[Amb]): Control[String] = {
    c = amb;
    pure("hello world")
}
AmbList { leaking }.run()
```

# Capability Safety

```
var c: Cap[Amb] = null
def leaking(implicit amb: Cap[Amb]): Control[String] = {
    c = amb;
    pure("hello world")
}
```

```
AmbList { leaking }.run()
{ flip()(c) }.run()
```

```
1. fish /Users/jonathan/AeroFS/work/github/algebraic-effects (fish)

hello 4

java.lang.RuntimeException: Prompt effekt.Capability$$anon$1@741e0037 not found on the stack.
  at scala.sys.package$.error(package.scala:27)
  at effekt.ReturnCont.splitAt(MetaCont.scala:24)
  at effekt.ReturnCont.splitAt(MetaCont.scala:19)
  at effekt.FramesCont.splitAt(MetaCont.scala:64)
  at effekt.Control$$anonfun$use$1.apply(Control.scala:121)
  at effekt.Control$$anonfun$use$1.apply(Control.scala:118)
  at effekt.Computation.apply(Control.scala:80)
  at effekt.Result$.trampoline(Result.scala:26)
  at effekt.Control$class.run(Control.scala:48)
  at effekt.Computation.run(Control.scala:78)
  at events.asyncPiping$suspended$$anonfun$interleave$1.apply(asyncPiping.scala:112)
  at events.asyncPiping$suspended$$anonfun$interleave$1.apply(asyncPiping.scala:108)
  at events.asyncPiping$Default$$anonfun$await$1$$anonfun$apply$4.apply(asyncPiping.scala:51)
  at events.asyncPiping$Default$$anonfun$await$1$$anonfun$apply$4.apply(asyncPiping.scala:51)
  at effekt.Control$$anonfun$use$1.apply(Control.scala:142)
  at effekt.Control$$anonfun$use$1.apply(Control.scala:118)
  at effekt.Computation.apply(Control.scala:80)
  at effekt.Result$.trampoline(Result.scala:26)
  at effekt.Control$class.run(Control.scala:48)
  at effekt.Computation.run(Control.scala:78)
  at events.asyncPiping$.runInterleaveUser(asyncPiping.scala:187)
... 43 elided

scala>
```

# Possible Solution

Make (capability) objects *second class* again

**Gentrification Gone too Far?**  
**Affordable 2nd-Class Values for Fun and (Co-)Effect**

Leo Osvald Grégory Escoffier Xilun Wu Lillian I. González-Alayón Dark Rompf  
Aurora University, USA; {losvald,gescorto,wu538,gonzal364,drmpf}@purdue.edu



**Abstract**  
First-class functions dramatically increase expressiveness at the expense of static guarantees. In ALGOL or PASCAL, functions could be passed as arguments but never escape their defining scope. Therefore, function arguments could serve as temporary access tokens or capabilities, enabling callers to perform some action, not only for the duration of the call. In modern languages, such programming patterns are no longer available.

The central thrust of this paper is to re-introduce second-class functions and other values alongside first-class entities in modern languages. We formalize second-class values with stack-bounded lifetimes as an extension to simply-typed  $\lambda$ -calculus, and for richer type systems such as  $F_\omega$ , and systems with path-dependent types. We generalize the binary first- vs second-class distinction to arbitrary privilege lattices, with the underlying type system as a special case. In this setting, abstract types naturally enable privilege polymorphism. We prove type soundness and lifetime properties in Coq.

We implement our system as an extension of Scala, and present several case studies. First, we modify the Scala Collections library and add privilege annotations to all higher-order functions. Privilege parametricity is key to retain the high degree of code-reuse between sequential and parallel as well as lazy and eager collections. Second, we use scoped capabilities to introduce a model of checked exceptions in the Scala library, with only few changes in the code. Third, we employ second-class capabilities for memory safety in a region-based off-heap memory library.

**Categories and Subject Descriptors:** D.3.3 [Programming Languages]: Languages—Concepts and Features

**Keywords:** first-class, second-class, types, effects, capabilities, object lifetimes

<sup>1</sup> Technically many languages still distinguish between, e.g., normal functions and closures, but not allow converting second-class first-class values via casts/operators, which effectively removes the distinction.

**A Study of Capability-Based Effect Systems**

Fengyun Liu

*A thesis submitted for the degree of Master of Computer Science  
at École Polytechnique Fédérale de Lausanne*

**Supervisors**

Martin Odersky Nada Amin Sandro Stucki  
Professor PhD Student PhD Student

School of Computer and Communication Sciences

**EPFL**  
ÉCOLE POLYTECHNIQUE  
FÉDÉRALE DE LAUSANNE

Lausanne, January 2016

# Second Class Values in Scala Escape

```
var c: Cap[Amb] = null
def leaking(implicit @local amb: Cap[Amb]): Control[String] = {
    c = amb;
    pure("hello world")
}
```

# Second Class Values in Scala Escape

```
var c: Cap[Amb] = null
def leaking(implicit @local amb: Cap[Amb]): Control[String] = {
    c = amb;
    pure("hello world")
}
```

**Error:** local value *amb* cannot be assigned to variable *c* since it would leave the scope of function *leaking*.

# Second Class Values in Scala Escape

```
var c: Cap[Amb] = null
def leaking(implicit @local amb: Cap[Amb]): Control[String] = {
    c = amb;
    pure("hello world")
}
```

**Error:** local value *amb* cannot be assigned to variable *c* since it would leave the scope of function *leaking*.

Restricts scope of capabilities so that they can be stack allocated.

# Second Class Values in Scala Escape

```
var c: Cap[Amb] = null
def leaking(implicit @local amb: Cap[Amb]): Control[String] = {
    c = amb;
    pure("hello world")
}
```

**Error:** local value *amb* cannot be assigned to variable *c* since it would leave the scope of function *leaking*.

Restricts scope of capabilities so that they can be stack allocated.  
This perfectly fits algebraic effects.

# 1 Effekt

## Solution Attempt 2

```
trait Person {  
    def greet(): Control[Unit]  
}
```

```
trait IOPerson extends Person {  
    implicit val console: Cap[Console]  
}
```

```
trait AlertPerson extends Person {  
    implicit val gui: Cap[GUI]  
}
```

# Solution Attempt 2

```
trait Person {  
    def greet(): Control[Unit]  
}
```

*Object lifetime < capability lifetime*

```
trait IOPerson extends Person {  
    implicit val console: Cap[Console]  
}
```

```
trait AlertPerson extends Person {  
    implicit val gui: Cap[GUI]  
}
```

# The Root of Evil

There is a simple connection:

- Algebraic effects and delimited continuations are all about the **stack**
- Object oriented programming is all about **heap allocated objects**

Conflicting requirements:

- capabilities should be stack allocated, objects don't
- but object lifetime should not be coupled to capability lifetime
- in particular, objects should be able to escape the handler scope losing the capabilities



# Part V

## Effectful

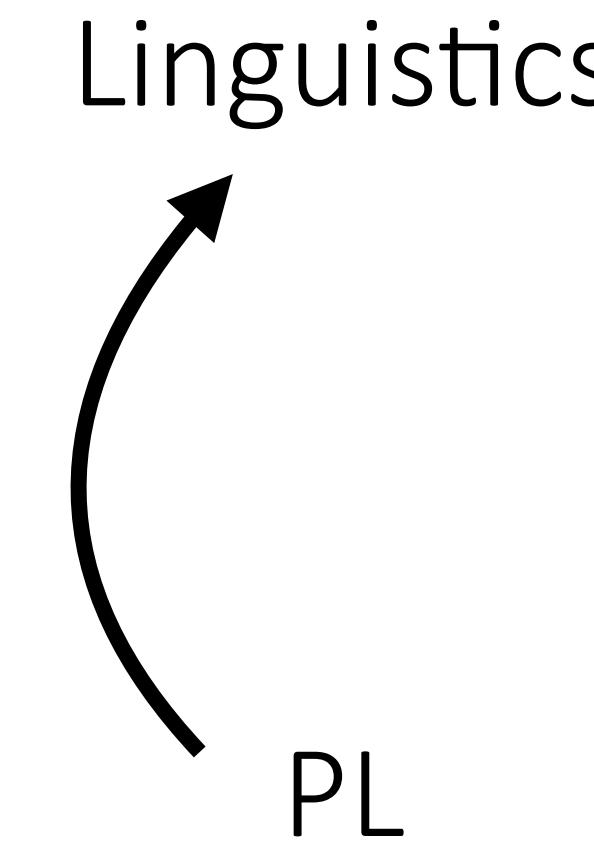
# Syntax

# Effectful Syntax

Linguistic phenomena like anaphora, scoping, quantification, implicature, focus and more can be modeled uniformly using algebraic effects.

## Algebraic Effects

Jiří Maršík and Maxime Amblard, 2016

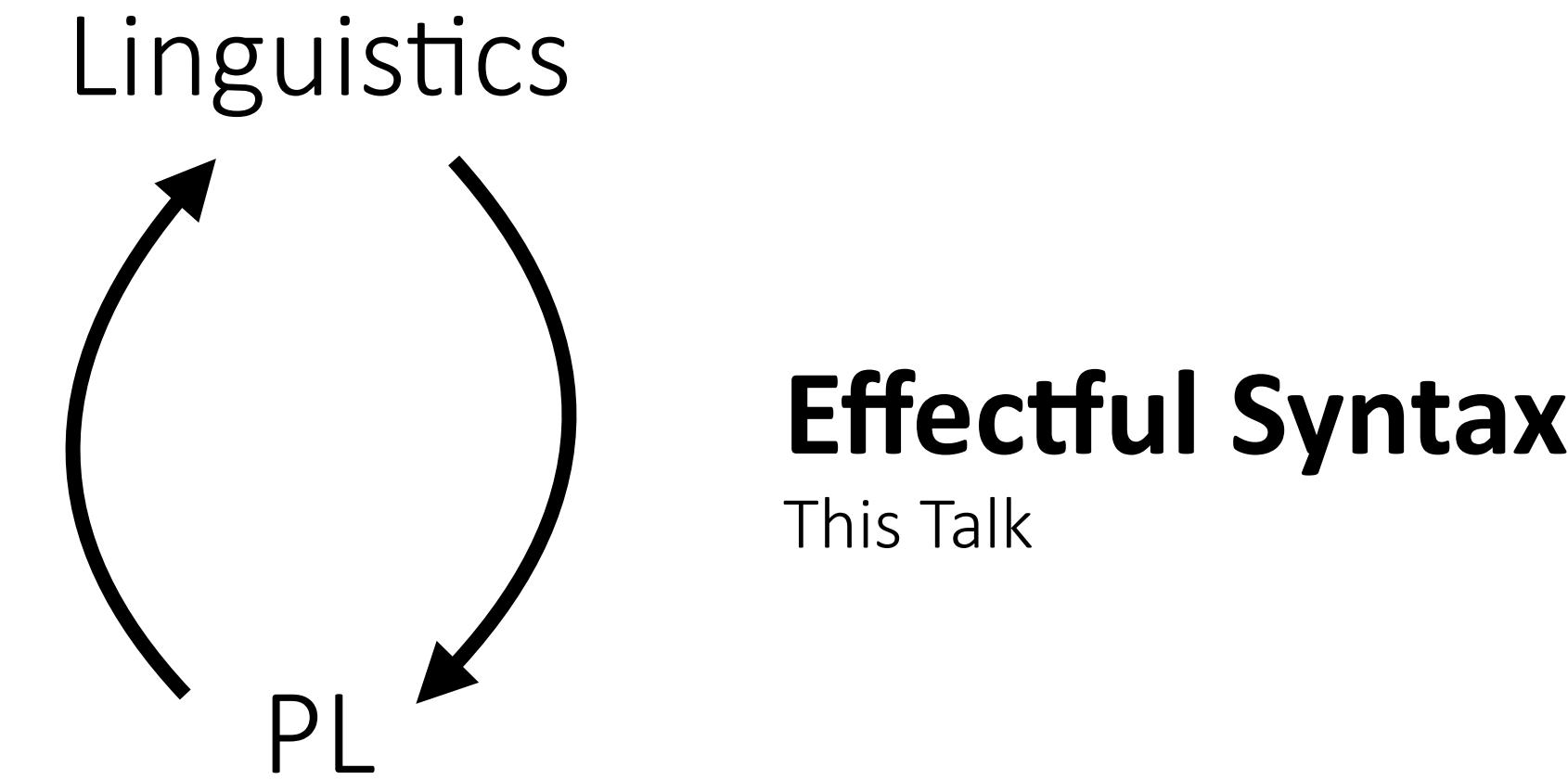


# Effectful Syntax

Linguistic phenomena like anaphora, scoping, quantification, implicature, focus and more can be modeled uniformly using algebraic effects.

## Algebraic Effects

Jiří Maršík and Maxime Amblard, 2016



- **Support linguistic phenomena in EDSLs using algebraic effects**
- **Use (algebraic) effects for AST construction**

# Example 1: The Speaker Effect

```
val s1: Sentence using Speaker = john said { mary loves me }
```

# Example 1: The Speaker Effect

Effect Signature

Groups effect operations in a type

```
val s1: Sentence using Speaker = john said { mary loves me }
```

Effect Operations

Semantics of the operations is left open

# Example 1: The Speaker Effect

Effect Signature

Groups effect operations in a type

```
val s1: Sentence using Speaker = john said { mary loves me }
```

|

|

Effect Operations

Semantics of the operations is left open

```
pete saidQuote { s1 }
```

# Example 1: The Speaker Effect

## Effect Signature

Groups effect operations in a type

```
val s1: Sentence using Speaker = john said { mary loves me }
```

|

|

## Effect Handlers

Provide semantics to effect operations

|

```
pete saidQuote { s1 }
```

## Effect Operations

Semantics of the operations is left open

# Example 1: The Speaker Effect

## Effect Signature

Groups effect operations in a type

```
val s1: Sentence using Speaker = john said { mary loves me }
```

## Effect Handlers

Provide semantics to effect operations

```
|  
pete saidQuote { s1 }
```

```
> Said(Pete, Said(John, Loves(Mary, Pete)))
```

## Effect Operations

Semantics of the operations is left open

## Example 2: The Scope Effect

```
val s2: Sentence using Scope = john saidQuote { every(woman) loves me }
```

## Example 2: The Scope Effect

```
val s2: Sentence using Scope = john saidQuote { every(woman) loves me }
```



Effect Operations

Semantics of the operations is left open

## Example 2: The Scope Effect

Effect Signature

Groups effect operations in a type

```
053   val s2: Sentence using Scope = john saidQuote { every(woman) loves me }
```

|

Effect Operations

Semantics of the operations is left open

## Example 2: The Scope Effect

Effect Signature

Groups effect operations in a type

```
053   val s2: Sentence using Scope = john saidQuote { every(woman) loves me }
```

|



Effect Operations

Semantics of the operations is left open

```
scoped { s2 }
```

## Example 2: The Scope Effect

053

Effect Handlers

Provide semantics to effect operations

```
scoped { s2 }
```

Effect Signature

Groups effect operations in a type

```
val s2: Sentence using Scope = john saidQuote { every(woman) loves me }
```

Effect Operations

Semantics of the operations is left open

## Example 2: The Scope Effect

053

Effect Handlers

Provide semantics to effect operations

```
|  
scoped { s2 }
```

```
> Forall(x => Implies(Woman(x), Said(John, Loves(x, John))))
```

Effect Signature

Groups effect operations in a type

```
val s2: Sentence using Scope = john saidQuote { every(woman) loves me }
```

Effect Operations

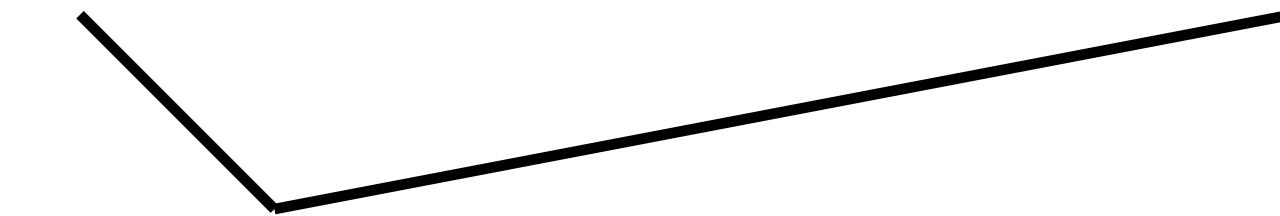
Semantics of the operations is left open

## Example 3: The Implicature Effect

```
val s3: Sentence using Speaker and Implicature =  
  mary loves { john whoIs { _ bestFriendOf me } }
```

## Example 3: The Implicature Effect

```
val s3: Sentence using Speaker and Implicature =  
  mary loves { john whoIs { _ bestFriendOf me } }
```



Effect Operations

Semantics of the operations is left open

## Example 3: The Implicature Effect

```
val s3: Sentence using Speaker and Implicature =  
    mary loves { john whoIs { _ bestFriendOf me } }  
  
pete saidQuote { accommodate { s3 } }
```

## Example 3: The Implicature Effect

```
val s3: Sentence using Speaker and Implicature =  
  mary loves { john whoIs { _ bestFriendOf me } }
```

```
pete saidQuote { accommodate { s3 } }
```

```
> Said(Pete,  
      And(BestFriendOf(John, Pete),  
           Loves(Mary, John)))
```

## Example 3: The Implicature Effect

```
val s3: Sentence using Speaker and Implicature =  
  mary loves { john whoIs { _ bestFriendOf me } }
```

```
pete saidQuote { accommodate { s3 } }
```

```
> Said(Pete,  
      And(BestFriendOf(John, Pete),  
           Loves(Mary, John)))
```

```
accommodate { pete saidQuote { s3 } }
```

## Example 3: The Implicature Effect

```
val s3: Sentence using Speaker and Implicature =  
  mary loves { john whoIs { _ bestFriendOf me } }
```

```
pete saidQuote { accommodate { s3 } }
```

```
> Said(Pete,  
      And(BestFriendOf(John, Pete),  
           Loves(Mary, John)))
```

```
accommodate { pete saidQuote { s3 } }
```

```
> And(Said(Pete, BestFriendOf(John, Pete)),  
      Said(Pete, Loves(Mary, John)))
```

# Effectful Syntax

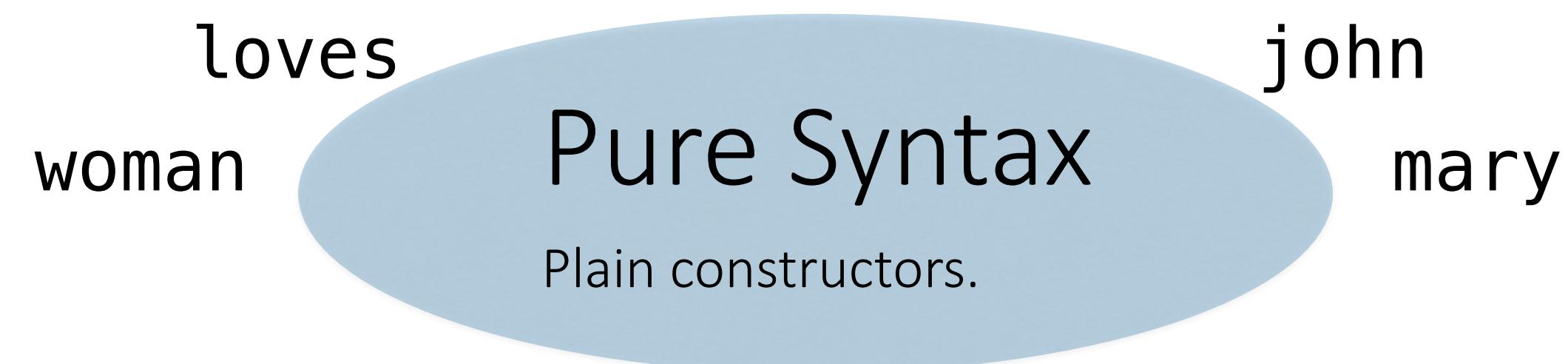
**Use (algebraic) effects for AST construction.**

More precisely, we propose to group syntax elements of DSLs into:

# Effectful Syntax

**Use (algebraic) effects for AST construction.**

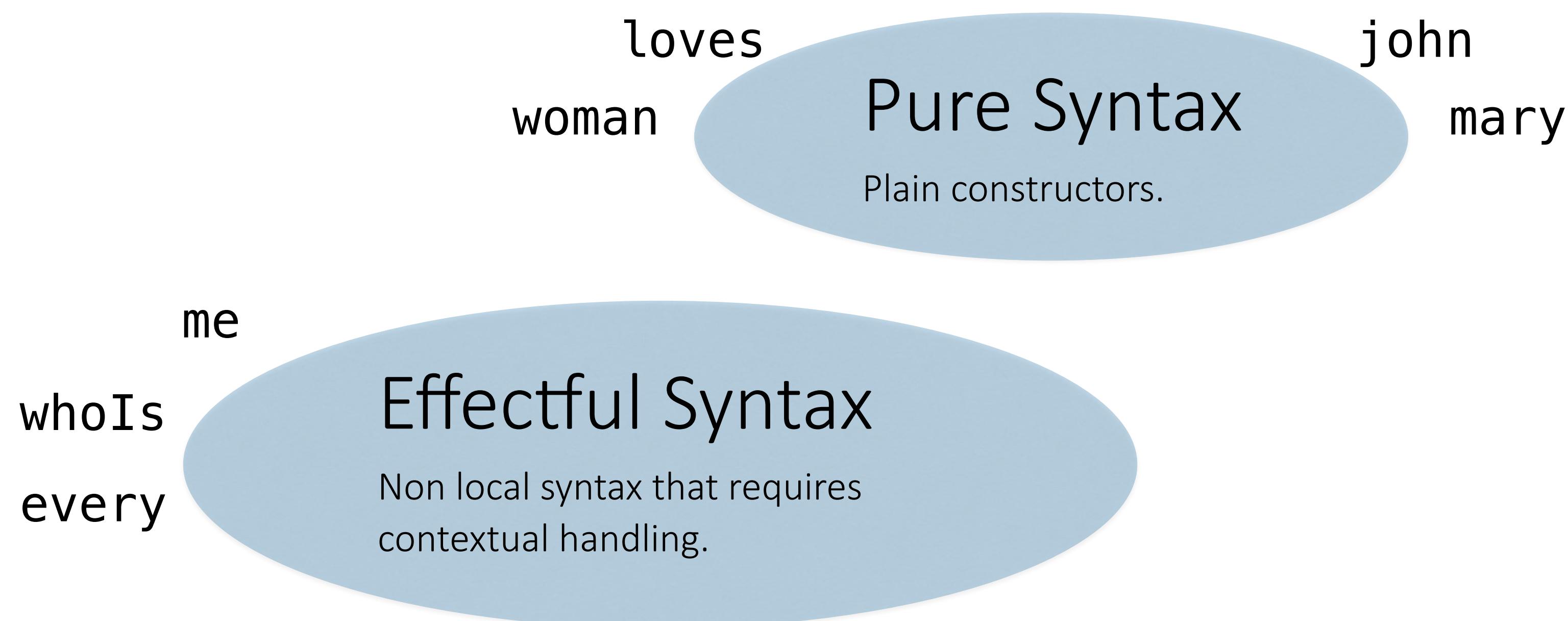
More precisely, we propose to group syntax elements of DSLs into:



# Effectful Syntax

**Use (algebraic) effects for AST construction.**

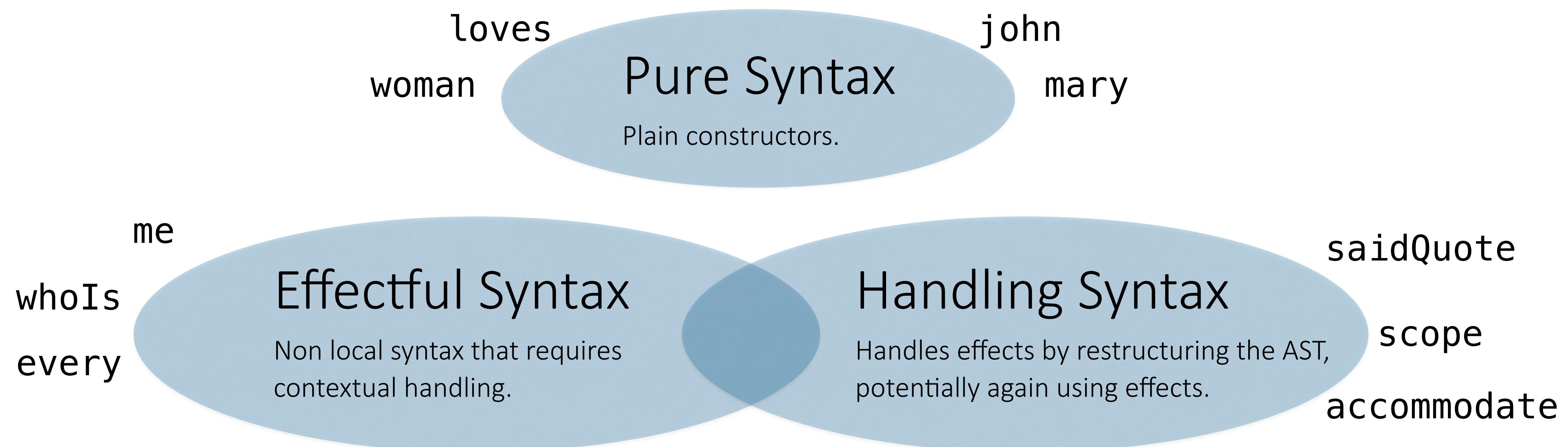
More precisely, we propose to group syntax elements of DSLs into:



# Effectful Syntax

**Use (algebraic) effects for AST construction.**

More precisely, we propose to group syntax elements of DSLs into:



# Properties

Effectful syntax based on algebraic effects is...

# Properties

Effectful syntax based on algebraic effects is...

## Modular

Linguistic phenomena  
can be encapsulated  
into reusable modules.

# Properties

Effectful syntax based on algebraic effects is...

## Modular

Linguistic phenomena  
can be encapsulated  
into reusable modules.

Effectful Syntax  
Handling Syntax

Pure Syntax

# Properties

Effectful syntax based on algebraic effects is...

## Modular

Linguistic phenomena  
can be encapsulated  
into reusable modules.

Effectful Syntax  
Handling Syntax

## Learnable

Separating linguistic  
phenomena from  
other domain concepts  
allows separate  
understanding

Pure Syntax

# Properties

Effectful syntax based on algebraic effects is...

## Modular

Linguistic phenomena can be encapsulated into reusable modules.

Effectful Syntax  
Handling Syntax

Pure Syntax

## Learnable

Separating linguistic phenomena from other domain concepts allows separate understanding

## Maintainable

Types precisely communicate usage of effectful syntax.